

GAME THEORY

Administrative stuff

1) Exercises

2) Billu "Oral exam" as "final project"

Summary

1) Game theory is about strategic decision-making
(Pareto \checkmark , Pareto \times)

2) Elements: Players, Actions, order of moves, payoffs, information

3) Actions: What each player can do

Strategies: Plan which action to pick given the information the player currently has

§ 2 Static games with complete information (SGCI)

Definition 2.1 (SGCI, "one-shot games", "normal-form games")

SGCI are those games for which

(1) players have all payoff-relevant information

(2) players move simultaneously.

(or in ignorance of what other players did.)

Elements of SGCI:

1) Players $N = \{1, \dots, n\}$

2) Action of player i is an element of

$$A^{(i)} = \{\alpha_1^{(i)}, \dots, \alpha_{j_i}^{(i)}\}$$

$$A = A^{(1)} \times \dots \times A^{(n)}$$

$$a \in A \quad a = (a_1, \dots, a_n)$$

3) Payoffs $\pi: A \rightarrow \mathbb{R}^n$

Remark 2.2 (Two-player games with finitely many actions)

$$N = \{1, 2\}$$

$$A^{(1)} = \{\alpha_1, \alpha_2, \dots, \alpha_m\}$$

$$A^{(2)} = \{\beta_1, \beta_2, \dots, \beta_k\}$$

$$\pi: A^{(1)} \times A^{(2)} \rightarrow \mathbb{R}^2$$

\leadsto usually written as a payoff matrix

	β_1	β_2	...	β_k
d_1	$\pi^{(1)}$	$\pi^{(2)}$		
d_2	"	"		
\vdots				
d_m	\square			$\pi_{mk}^{(1)}, \pi_{mk}^{(2)}$
	$\pi_{m1}^{(1)}, \pi_{m1}^{(2)}$			

Similar way how to represent 3-player games
 ↪ Exercises

Example 2.3 (Prisoner's dilemma)

Two players, they can either confess or remain silent

Payoffs $\hat{=}$ saved prison time

	silent	confess		S	C
silent	3, 3	0, 4	S	3	0
confess	4, 0	1, 1		C	4

This game has the following properties

1) $A^{(1)} = A^{(2)}$

2) $\pi^{(1)}(a, a') = \pi^{(2)}(a', a)$

2-player games with this property are called "symmetric"

n-player games? ↪ Exercises

Definition 2.4 (Strategies & SGC1)

1) Let $T = (N, A, \pi)$ be a SGC1.

$$\text{let } A^{(i)} = \{a_1^{(i)}, \dots, a_m^{(i)}\}$$

Then a strategy for player i is a probability distribution over the set of actions

$$\sigma^{(i)} = (\sigma_1^{(i)}, \sigma_2^{(i)}, \dots, \sigma_m^{(i)}) \quad \sigma_j \geq 0 \quad \sum_j \sigma_j = 1$$

Set of all strategies $\Sigma^{(i)}$

Set of "strategy profiles" : $\Sigma = \Sigma^{(1)} \times \Sigma^{(2)} \times \dots \times \Sigma$

2) A strategy is called pure

if there is a j such that $\sigma_j^{(i)} = 1$, $\sigma_k^{(i)} = 0$
for all $k \neq j$.

If I would like to highlight that a strategy is pure

then I write $s = (0, \dots, 0, 1, 0, \dots, 0)$

Set of all pure strategies of player i is $S^{(i)}$

Example 2.5 (Prisoner's dilemma)

Strategy for player 1 : $(1/2, 1/2)$

Probability to remain
fidel

Probability to confess

Pure strategies $(1, 0)$

(0, 1)

Remark 2.6 (Payoffs for mixed strategies)

(*) If all players use pure strategies

the resulting payoff is $\pi(\omega^{(1)}, \omega^{(2)}, \dots, \omega^{(n)})$

(*) I want $\pi: \Sigma \rightarrow \mathbb{R}^n$

(*) Idea: Each player randomizes independently
Strategies are $\sigma^{(i)} = (\sigma_j^{(i)})$ $\sigma := (\sigma^{(1)}, \sigma^{(2)}, \dots, \sigma^{(n)}) \in \Sigma$

$$\pi(\sigma) := \sum_{j_1, \dots, j_n} \sigma_{j_1}^{(1)} \cdot \sigma_{j_2}^{(2)} \cdot \dots \cdot \sigma_{j_n}^{(n)} \pi(\omega_{j_1}^{(1)}, \omega_{j_2}^{(2)}, \dots, \omega_{j_n}^{(n)})$$

Example 2.7 (Prisoner's dilemma)

$A^{(i)} = \{ \text{Silent}, \text{Confess} \}$

$$\sigma^{(1)} = (0.6, 0.4)$$

$$\sigma^{(2)} = (0.3, 0.7)$$

	S	C
S	3	0
C	4	1

Possible outcomes	Probability	Payoff of
(Silent, Silent)	$0.6 \cdot 0.3$	3
(Silent, Confess)	$0.6 \cdot 0.7$	0
(Confess, Silent)	$0.4 \cdot 0.3$	4
(Confess, Confess)	$0.4 \cdot 0.7$	1

$$\begin{aligned} \pi^{(1)}(\sigma^{(1)}, \sigma^{(2)}) &= 0.6 \cdot 0.3 \cdot 3 \\ &\quad + 0.6 \cdot 0.7 \cdot 0 \\ &\quad + 0.4 \cdot 0.3 \cdot 4 \\ &\quad + 0.4 \cdot 0.7 \cdot 1 \end{aligned}$$

		0.3	0.7
		S	C
0.6	S	3	0
0.4	C	4	1

$\underbrace{\hspace{10em}}_P$
 $\sigma^{(1)} \cdot P \sigma^{(2)T}$

Notation 2.8

If I want to speak of player i 's strategy specifically

I sometimes write $(\sigma^{(i)}, \sigma^{(-i)}) = (\sigma^{(1)}, \sigma^{(2)}, \dots, \sigma^{(n)})$

Similarly $\pi^{(i)}(\sigma^{(1)}, \dots, \sigma^{(n)}) = \pi^{(i)}(\sigma^{(i)}, \sigma^{(-i)})$

Remark 2.9 (Solution concepts)

So far: introduce machinery to describe a game.

Question: What does it mean to solve a game?

[What would rational players do?]

Solution concepts

- (1) Elimination of dominated strategies
- (2) Nash equilibria

§ 2.1 Iterated elimination of dominated strategies

Example 2.10 (Prisoner's dilemma)

	↓		
	Silent	Confess	
→ Silent	3,3	0,4	
→ Confess	4,0	1,1	

For rational players, (Confess, Confess) seems to be the only logical outcome.

Definition 2.11 (Dominated strategies)

(1) A ~~pre~~ strategy $s^{(i)}$ is strictly dominated

Suppose $\pi(s^{(i)}, s^{(-i)}) < \pi(\tilde{s}^{(i)}, s^{(-i)})$

↓ ↘ to more...

where does this live
is there some restriction

if there is some $\tilde{\sigma}^{(i)} \in \Sigma^i$ such that
for all $s^{(-i)} \in S^{(-i)}$:

$$\pi^{(i)}(s^{(i)}, s^{(-i)}) < \pi^{(i)}(\tilde{\sigma}^{(i)}, s^{(-i)})$$

(2) A pure strategy $s^{(i)}$ is weakly dominated
if

$$\pi^{(i)}(s^{(i)}, s^{(-i)}) \leq \pi^{(i)}(\tilde{\sigma}^{(i)}, s^{(-i)}) \quad \forall s^{(-i)}$$

$$\exists \tilde{s}^{(-i)} : \pi^{(i)}(s^{(i)}, \tilde{s}^{(-i)}) < \pi^{(i)}(\tilde{\sigma}^{(i)}, \tilde{s}^{(-i)})$$

$$\tilde{\sigma}^{(i)} := s^{(i)}$$

Remark 2.12 (on dominated strategies)

1) In prisoner's dilemma, "silent" is strictly dominated,

2) Why do we explicitly allow for domination by mixed strategies.

These are games where no strategy is dominated if you only allow dominated

by pure strategies, but still dominated
by mixed strategies is possible.

	L	R
T	3,0	0,0
M	1,1	1,1
D	0,0	3,0

Strategy M is dominated

↪ Exercises

3) There is no added advantage
of replacing $s^{(-i)}$ by $\sigma^{(-i)}$

If statement is true for all $s^{(-i)}$
it is also true for all $\sigma^{(-i)}$.

Example 2-13 (A prisoner's dilemma with remorse)

	⁽²⁾ Silent	Confess	
Silent	3,3	0,0	(1)
Confess	4,0	<u>1,1</u>	